

Graph algorithms and Breadth-first search

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Outline

- ▶ What is graph search?
- ▶ What is breadth-first search (BFS)?

Graph algorithms

We said in previous lectures that what graph representation is best depends on what sort of questions we want to ask about a graph.

Questions we might ask are things like:

- ▶ What is the sum of all the vertices in a graph G , where each vertex holds an `int`?
- ▶ What is the shortest weighted path from one vertex (say, a vertex representing Perth airport) to another (say, a vertex representing Chongqing Jiangbei International Airport)?
- ▶ What is the least costly way of connecting vertices representing electric power stations and buildings?

Graph algorithms

Questions like

- ▶ What is the sum of all the vertices in a graph G , where each vertex holds an `int`?

are examples of *graph traversal* – they ask us to visit every vertex in the graph.

Two possible ways of performing graph traversal we could consider are:

- ▶ breadth-first search (BFS):
 - ▶ From an initial vertex, imagine starting *multiple* paths to neighbours, and advance each path one step at a time
- ▶ depth-first search (DFS):
 - ▶ Put the edges leading out of a starting vertex in some order, left to right, and continue exploring the left-most path out of each vertex until you reach an end to the path. Then start exploring one-after-the-left, then the second after the left, and so on (“deepest” first)

Depth-first search (DFS)

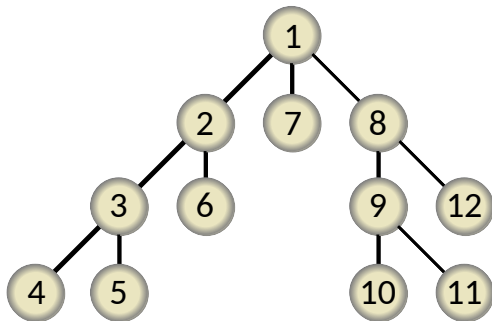
We have already encountered these terms (BFS and DFS) when looking at *trees*.

We said that for trees, depth-first search results from doing a “pre-order” traversal of a tree.

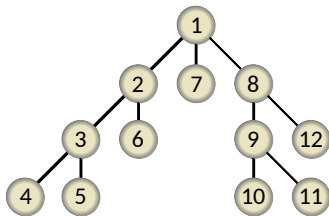
- ▶ Recall that the mnemonic for pre-order binary tree traversal is “NLR” (visit a *Node*, then its *Left* child, then the *Right*, recursively).

Depth-first search (DFS)

The following diagram shows the order nodes would be visited in a graph when doing a depth-first search.



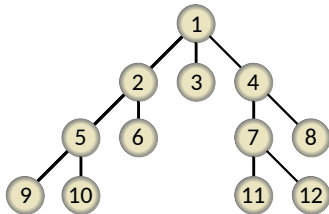
Depth-first search (DFS)



This graph happens to be a *tree*; trees are just a subset of graphs, where nodes in the tree are vertices in the graph, and an edge exists between any child node in the tree and its parent.

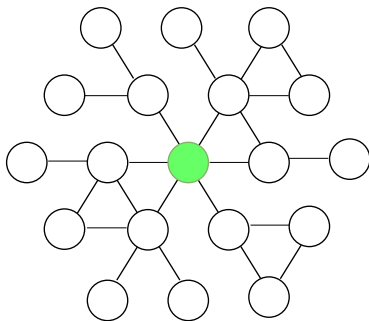
Breadth-first search (BFS)

Recall that for a tree, breadth-first search visits nodes “level by level”; we can do the same for graphs.



Breadth-first search

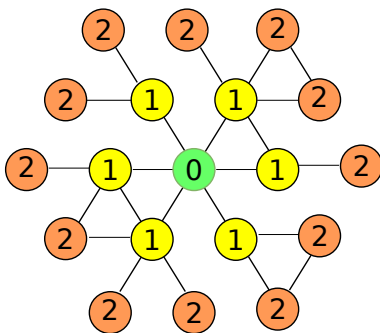
We can apply breadth-first search to graphs which are *not* trees, as well, like the following:



- We start with an initial node n , then we visit nodes which are one link distant from n , then 2, and so on.

Breadth-first search

The order we visit the nodes would be as follows



(where nodes with the same number are visited in some arbitrary order).

Previously, we only explained breadth-first search informally; we will now explain an algorithm for doing breadth-first search.

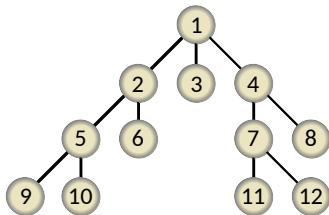
Breadth-first search (BFS)

This algorithm starts from a given vertex v and explores a graph G in a *breadth-first* manner.

We shall see also that as it does so, it constructs what we call a *spanning tree* for graph G , called the breadth-first tree.

Breadth-First Search Algorithm (BFS)

- ▶ The algorithm uses a **queue** as its main data structure.
- ▶ The basic idea is simple:
 - ▶ Starting with a queue containing just one node, consider the neighbours of each node.
 - ▶ Visit any new vertices from the neighbours, until there are no more vertices to be added.



Implementing BFS

In order to implement BFS we maintain three data structures:

1. **Q**, the queue of vertices to be processed
2. An array **colour**, with values white (0), grey (1) or black (2) for each vertex.
3. An array **p**, which where **p**(v) contains the immediate parent of v in the spanning tree. That is the edges of the spanning tree are given by (v, **p**(v))

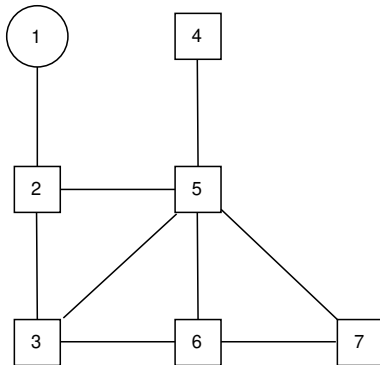
BFS Algorithm

- put v into the queue Q
- $\text{colour}[v] = \text{white}$ for all vertices v ,
- $p(v)$ is undefined.
- while Q is not empty:
 - dequeue w from the head of Q
 - for each vertex x adjacent to w :
 - if $\text{colour}[x]$ is white:
 - $p(x) = w$
 - $\text{colour}[x] = \text{grey}$
 - enqueue x onto the tail of Q
 - $\text{colour}[w] = \text{black}$

BFS Algorithm

At the end of the BFS search, every node will have the colour black and the parent array p will contain details of a BFS spanning tree.

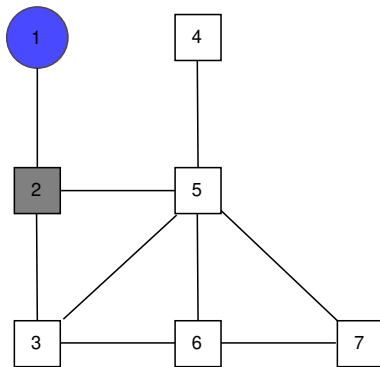
Let us examine how the BFS algorithm works on the following graph:



Initially, our queue Q contains 1.

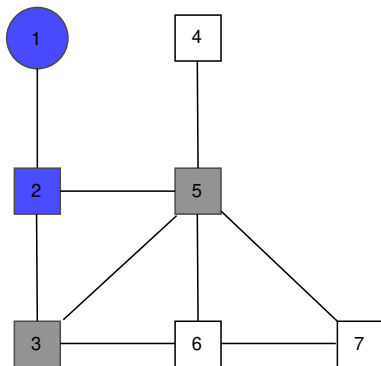
We will ignore the use of the p array for the moment – but basically, each time we add something to the queue, we use p to store where the search came *from* when moving to some node n .

BFS step 1



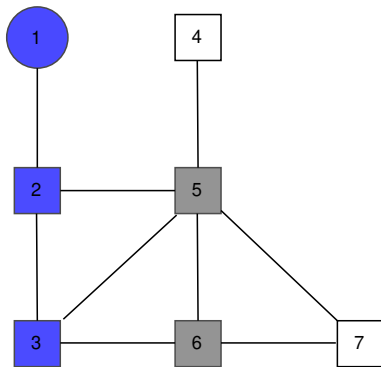
- ▶ $Q = [1]$
- ▶ Dequeue 1 from the queue Q
- ▶ Colour 2 grey and add it to the queue
- ▶ Colour 1 black

BFS step 2



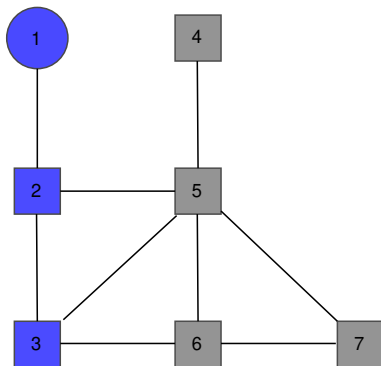
- ▶ $Q = [2]$
- ▶ Dequeue 2 from the queue Q
- ▶ Colour 3 grey and add it to the queue
- ▶ Colour 5 grey and add it to the queue
- ▶ Colour 2 black

BFS step 3



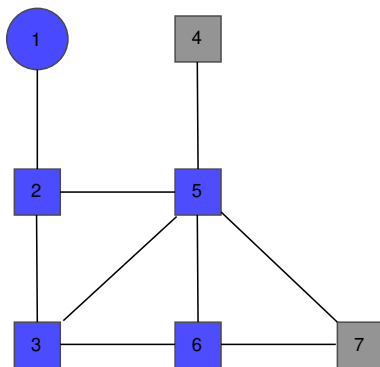
- ▶ $Q = [3, 5]$
- ▶ Dequeue 3 from the queue Q
- ▶ Colour 6 grey and add it to the queue
- ▶ Colour 3 black

BFS step 4



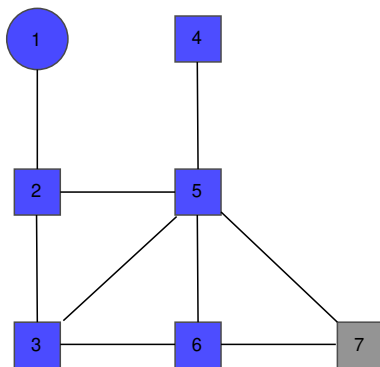
- ▶ $Q = [5, 6]$
- ▶ Dequeue 5 from the queue Q
- ▶ Colour 4 grey and add it to the queue
- ▶ Colour 7 grey and add it to the queue
- ▶ Colour 5 black

BFS step 5



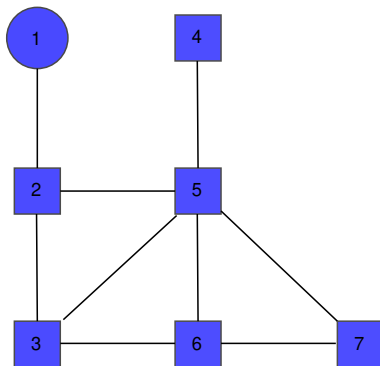
- ▶ $Q = [6, 4, 7]$
- ▶ Dequeue 6 from the queue Q
- ▶ Colour 6 black

BFS step 6



- ▶ $Q = [4, 7]$
- ▶ Dequeue 4 from the queue Q
- ▶ Colour 4 black

BFS step 7



- ▶ $Q = [7]$
- ▶ Dequeue 7 from the queue Q
- ▶ Colour 7 black

BFS Java code

The code for the BFS algorithm is provided in `BFS.java` in your Java code bundle.

BFS and “shortest path” problems

Breadth-first search also provides us with a way of answering questions like

- ▶ Given two vertices A and B in a graph G , what is the shortest path between them?

BFS and “shortest path” problems

- ▶ Given two vertices A and B in a graph G , what is the shortest path between them?

To answer this question, we just do breadth-first search *starting* from vertex A .

And each time we visit a vertex, we check whether we've found vertex B ; if we have, we stop. And the p array, which stores the “path back” to node A , is the shortest route from A to B .